A characterisation of functions computable in polynomial time and space over the reals with discrete ordinary differential equations

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## Introduction

- providing algebraic characterisations of polynomial complexity classes, without:
  - explicit bounds (Cobham)
  - several types of arguments (Bellantoni and Cook)
  - tiering (Leivant and Marion)...
- We are in the framework of implicit complexity

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- ... but here we are interested in functions over the real numbers

- Computable real number
- Computable function over the reals
- Discrete derivation :  $f : \mathbb{N} \times \mathbb{R} \mapsto \mathbb{R}$ ,

$$f'(x,y) = \frac{\partial f}{\partial x}(x,y) = f(x+1,y) - f(x,y).$$

Previous work: characterising computable reals and sequences (MCU22)

### Definition (Length ODE)

A function **f** is "length-ODE" definable (from u, g and h) if it is a solution of

$$f(0,\mathbf{y}) = \mathbf{g}(\mathbf{y})$$
 and  $\frac{\partial \mathbf{f}(x,\mathbf{y})}{\partial \ell} = \mathbf{u}(\mathbf{f}(x,\mathbf{y}),\mathbf{h}(x,\mathbf{y}),x,\mathbf{y}).$  (1)

Formal synonym for right-hand side of (1):  $\mathbf{f}(x+1,\mathbf{y}) = \mathbf{f}(x,\mathbf{y}) + (\ell(x+1) - \ell(x)) \cdot \mathbf{u}(\mathbf{f}(x,\mathbf{y}),\mathbf{h}(x,\mathbf{y}),x,\mathbf{y})$  A derivation with respect to the length corresponds to change of variable in the ODE. Variation when:

 $\ell(x+1)-\ell(x)\neq 0$ 

Inspired by:

$$\frac{\delta f(x,\mathbf{y})}{\delta x} = \frac{\delta \ell(x)}{\delta x} \cdot \frac{\delta f(x,\mathbf{y})}{\delta \ell(x)}.$$

$$f(0) = 2$$
  
$$\frac{\partial f}{\partial \ell}(x) = f(x) \cdot f(x) - f(x)$$

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Unique solution :  $f(x) = 2^{2^{\ell(x)}}$ 

**Definition**  $f : \mathbb{R} \to \mathbb{R}$  is an effective limit of  $\overline{f} : \mathbb{R} \times \mathbb{N} \to \mathbb{R}$  if  $|f(x) - \overline{f}(x, 2^n)| \le 2^{-n}$ 

Key observation :  $\overline{f}$  computable in polynomial time  $\Rightarrow f$  computable in polynomial time.

We consider:

$$\overline{\mathbb{LDL}^{\bullet}} = [\mathbf{0}, \mathbf{1}, \pi_i^k, \ell(x), +, -, \times, \overline{\mathrm{cond}}(x), \frac{x}{2};$$

composition, linear length ODE, effective limit]

Theorem (MCU22)  $\mathsf{FPTIME} \cap \mathbb{R}^{\mathbb{N}} = \overline{\mathbb{LDL}^{\bullet}} \cap \mathbb{R}^{\mathbb{N}}$ 



# Characterising computable functions over the reals



We consider:

$$\overline{\mathbb{LDL}^{\circ}} = [\mathbf{0}, \mathbf{1}, \pi_i^k, \ell(x), +, -, \tanh, \frac{x}{2}, \frac{x}{3};$$

composition, linear length ODE, effective limit]

Theorem

$$\overline{\mathbb{L}\mathbb{D}\mathbb{L}^{\circ}}\cap\mathbb{R}^{\mathbb{R}}\ =\mathsf{FPTIME}\cap\mathbb{R}^{\mathbb{R}}$$

#### Definition (Robust linear ODE)

A bounded function **f** is a robustly linear ODE definable (with **u** essentially linear in  $\mathbf{f}(x, \mathbf{y})$ , **g** and **h**) if:

1. it is a solution of

$$\mathbf{f}(0,\mathbf{y}) = \mathbf{g}(\mathbf{y})$$
 and  $\frac{\partial \mathbf{f}(x,\mathbf{y})}{\partial x} = \mathbf{u}(\mathbf{f}(x,\mathbf{y}),\mathbf{h}(x,\mathbf{y}),x,\mathbf{y}),$ 

2. the ODE is (polynomially) numerically stable.

We consider:

$$\overline{\mathbb{RLD}^{\circ}} = [\mathbf{0}, \mathbf{1}, \pi_i^k, \ell(x), +, -, \tanh, \frac{x}{2}, \frac{x}{3};$$

composition, robust linear ODE, effective limit]

Theorem

 $\overline{\mathbb{RLD}^{\circ}} \cap \mathbb{R}^{\mathbb{R}} = \mathsf{FPSPACE} \cap \mathbb{R}^{\mathbb{R}}$ 

## Conclusion

- We now have "natural" algebraic characterisations for functions over the reals computable in FPTIME and FPSPACE.
- Further works:
  - Going to continuous ODEs
  - Cover other complexity classes
  - See what we can say for more general functions